

1 SUMMARY

Calculates **the minimax solution** of a system of m **linear algebraic equations** in n unknowns, $m \geq n$, where the maximum element of the solution is **subject to a simple bound**. Given equations

$$\sum_{j=1}^n a_{ij}x_j + b_i = 0 \quad i=1,2,\dots,m \quad m \geq n$$

find the solution $x_j, j=1,2,\dots,n$ such that

$$\max_i \left\{ \left| \sum_{j=1}^n a_{ij}x_j + b_i \right| \right\}$$

is minimized subject to the bounds $|x_j| \leq g, j=1,2,\dots,n$.

A variation of the 'exchange algorithm' is used that incorporates a technique for reducing the number of iterations, and which will also provide a defined solution even when the matrix $\mathbf{A} = \{a_{ij}\}_{n \times m}$ is rank deficient.

Described in K.Madsen and M.J.D.Powell, Harwell report R.7954 (1975).

ATTRIBUTES — **Version:** 1.0.0. **Types:** MA19A; MA19AD. **Original date:** March 1974. **Origin:** K.Madsen, Copenhagen.

2 HOW TO USE THE PACKAGE

2.1 The argument lists

The single precision version

```
CALL MA19A(N,M,A,IA,B,G,EPS,X,RES,IREF)
```

The double precision version

```
CALL MA19AD(N,M,A,IA,B,G,EPS,X,RES,IREF)
```

- N** is an INTEGER variable and must be set by the user to n , the number of unknowns. **Restriction:** $n > 0$.
- M** is an INTEGER variable and must be set to m , the number of linear expressions under consideration. **Restriction:** $m > 0$.
- A** is a REAL (DOUBLE PRECISION in the D version) two-dimensional array which the user must set to the coefficients a_{ij} , i.e. $A(i, j) = a_{ij}, i=1,2,\dots,m, j=1,2,\dots,n$. It is not changed by the subroutine.
- IA** is an INTEGER variable and must be set to the first dimension of the array A.
- B** is a REAL (DOUBLE PRECISION in the D version) array which the user must set to the constants b_i , i.e. $B(i) = b_i, i=1,2,\dots,m$. It is not changed by the subroutine.
- G** is a REAL (DOUBLE PRECISION in the D version) variable which must be set to the value of the bound g on the unknowns x_j . **Restriction:** $g \geq 0$.
- EPS** is a REAL (DOUBLE PRECISION in the D version) variable which controls the accuracy of the solution, i.e. the solution will satisfy equation (3) in section 4 with $\epsilon = \text{EPS}$. It has to be set to a non-negative value by the user, and normally $\text{EPS} = 0$ will be a good value.
- X** is a REAL (DOUBLE PRECISION in the D version) array in which the subroutine will return the solution. The array should be of length at least n .

RES is a REAL (DOUBLE PRECISION in the D version) array which is used for working space. Its length must be at least $(n+1)(n+5)+m$. On exit it will contain the values

$$\text{RES}(i) = \sum_{j=1}^n a_{ij}x_j + b_i, \quad i=1,2,\dots,m.$$

Further, $\text{RES}(m+1)$ will contain the value of h in equation (1) in section 4, and $\text{RES}(m+2)$ will contain an upper bound for the expressions on the left-hand side of equation (3) in section 4.

IREF is an INTEGER array which is used for working space. Its length must be at least $4(n+1)+m$. On exit $|\text{IREF}(i)|$, $i=1,2,\dots,(n+1)$, will contain information about which equations belong to the sets I_D and I_B defined in section 4.

$|\text{IREF}(i)| > m$ will mean that the index $|\text{IREF}(i)| - m$ is in the set I_B . Otherwise $|\text{IREF}(i)|$ belongs to the set I_D . The sign of $\text{IREF}(i)$ is the sign of s_i in equations (1) and (2) in section 4. Further $\text{IREF}(n+2)$ will give the number of elements in I_B , and $\text{IREF}(n+3)$ will give the number of iterations used by the method.

3 GENERAL INFORMATION

Use of common: None.

Workspace: Provided by the user, see arguments RES and IREF.

Other routines called directly: None.

Input/output: None.

Restrictions: $n \geq 1$, $m \geq 1$, $g \geq 0$, $\varepsilon \geq 0$.

4 METHOD

The method will be described in a forthcoming Harwell report. The solution to the problem will be the solution to $(n+1)$ linear equations:

$$\sum_{j=1}^n a_{ij}x_j + b_i = s_i h, \quad \text{for } i \in I_D \quad (1)$$

$$x_i = s_i g_1 \quad \text{for } i \in I_B \quad (1)$$

where $h \geq 0$, and $g_1 = g$ unless $h=0$ when $g_1 \leq g$. Further $s_i = \pm 1$, and I_B contains n_b indices, I_D contains $(n+1-n_b)$ indices, $0 \leq n_b \leq n$. At the solution the inequalities

$$\left| \sum_{j=1}^n a_{ij}x_j + b_i \right| \leq h + \varepsilon \quad (3)$$

and

$$|x_k| \leq g$$

will be satisfied for all relevant values of i and k .

Note that if the bounds instead of being of the form $|x_j| \leq g$, $j=1,2,\dots,n$ are

$$|x_j| \leq g_j, \quad j=1,2,\dots,n$$

the problem is easily re-written in the first form by scaling the coefficients a_{ij} .