## 1 SUMMARY

To compute the modified Bessel functions Ibessnu and $K v$ for a range of orders $v=0,1,2, \ldots, n$, at the same argument $x$.

ATTRIBUTES - Version: 1.0.0. Types: FF07A; FF07AD. Calls: FD05A. Original date: April 1983. Origin: A. R. Curtis, Harwell.

## 2 HOW TO USE THE PACKAGE

### 2.1 Argument list and calling sequence

The single precision version
CALL FF07A (X, N, A, B, NAB)
The double precision version

```
CALL FF07AD (X,N,A,B,NAB)
```

$\mathrm{X} \quad$ is a REAL (DOUBLE PRECISION in the D version) variable which must be set by the user to the value of $\pm x$ (the subroutine uses the absolute value of X$)$. A zero value is not allowed when $K_{v}(x)$ is requested, i.e. when argument $\mathrm{N}<0$. This argument is not altered by the subroutine.
$\mathrm{N} \quad$ is an INTEGER variable which must be set by the user to $\pm n$, where $n$ is the highest order required. If $N<0$, both $I_{v}(x)$ and $K_{v}(x)$ are computed; if $\mathrm{N}>0$, only $I_{v}(x)$ is computed. A zero value is not allowed. This argument is not altered by the subroutine.

A is a REAL (DOUBLE PRECISION in the D version) array of dimension NAB in which the subroutine will return the values of $I_{v}(x), v=0,1,2, \ldots, n$, i.e. $I_{0}(x)$ in A (1) , $I_{1}(x)$ in A (2) up to $I_{n}(x)$ in $\mathrm{A}(\mathrm{n}+1)$.
B is a REAL (DOUBLE PRECISION in the D version) array of dimension NAB in which the subroutine will return the values of $K_{v}(x), v=0,1,2, \ldots, n$, i.e. $K_{0}(x)$ in $\mathrm{B}(1), K_{1}(x)$ in $\mathrm{B}(2)$ up to $K_{n}(x)$ in $\mathrm{B}(\mathrm{n}+1)$. If $\mathrm{N}>0$ the array B is not altered.

NAB is an INTEGER variable which must be set by the user to the dimension of the arrays A and B; it must be at least $n+1$. This argument is not altered by the subroutine.

## 3 GENERAL INFORMATION

Use of common: none.
Workspace: None.
Other routines called directly: calls DEXP, DSQRT and DLOG.
Input/output: Error warning and diagnostic messages on unit 6 .

## Restrictions:

$n>0$;
$x>0$ if $K_{v}(x)$ required,
else $x \geq 0$.

## 4 METHOD

### 4.1 Method

The method is based on the recurrence relation (9.6.26 in Abramowitz and Stegun ${ }^{[1]}$ )

$$
\begin{equation*}
u_{v+1}-u_{v-1}=\frac{2 v}{x} u_{v} \tag{1}
\end{equation*}
$$

whose general solution is

$$
\begin{equation*}
u_{v}=a I_{v}(x)+b(-1)^{v} K_{v}(x) \tag{2}
\end{equation*}
$$

where $a$ and $b$ are constants.
To compute $I_{v}(x)$, the recurrence relation is solved for $u_{v-1}$ and used downwards from $v=m>n$, chosen so that the second term in (2) is negligible compared with the first for $v \leq n$. At the same time, the sum

$$
\begin{equation*}
s=u_{0}+2 \sum_{v=1}^{m} u_{v} \approx a e^{x} \tag{3}
\end{equation*}
$$

is computed, and the normalizing factor $a^{-1}=s^{-1} e^{x}$ is finally applied to the stored values in the array A.
To compute $K_{v}(x)$ (if required), the recurrence relation is solved for $u_{v+1}$ and used upwards from accurate values of $K_{0}(x)$ and $K_{1}(x)$. The first is computed from new highly accurate Chebyshev series, and $K_{1}(x)$ is then computed as (9.6.15 ${ }^{[1]}$ )

$$
\begin{equation*}
K_{1}(x)=\frac{1 / x-I_{1}(x) K_{0}(x)}{I_{0}(x)} \tag{4}
\end{equation*}
$$

which does not suffer serious loss of accuracy through cancellation.

### 4.2 Accuracy and timing

The computation of $I_{0}(x)$ is good almost to full computer accuracy. The values of $I_{v}(x)$ in Table 9.11 in Abramowitz and Stegun ${ }^{[1]}$ are reproduced exactly (apart from those which underflow). It may be necessary to take precautions against overflow. For small $x$ and large $n$, underflows will occur and zero values will be returned; a call could be made to a library subroutine to mask off underflow interrupts.

The accuracy of $K_{v}(x)$ is also about 15 significant figures using the new Chebyshev series for $K_{0}(x)$. All $K_{v}(x)$ values in Table $9.11^{[1]}$ for $x \leq 10$ are reproduced (except for a few discrepancies by one in the last digit of the table values, and for values which would overflow); however, for $x=50$ and $x=100$, many errors in the last digit of the table values where found, and the accuracy of FFO7AD was confirmed independently by using the asymptotic series 9.7.2 $2^{[1]}$. For small $x$ and large $n$, overflows would occur on the IBM computer; these are avoided, and the largest value represented on the machine is returned instead.

A 4-byte arithmetic version on some computers is not recommended because of possible accumulation of rounding errors; since the result arrays A and B are used as work-space, this consideration prevents internal 8-byte arithmetic computation being used to give 4-byte precision results. If FF07AD is used on other computers, the constants used in accuracy and overflow tests should be reconsidered.

Execution time depends more or less linearly on $n$, and also somewhat on $x$. For moderate $x$ values, it varies on an IBM/3081 from about $75 \mu \mathrm{sec}$ at $n=10$ to about $650 \mu \mathrm{sec}$ at $n=100$ for $I_{v}(x)$ only; if $K_{v}(x)$ is also computed, typical times are $125 \mu \mathrm{sec}$ and $800 \mu \mathrm{sec}$. Values of $x$ in the range 10 to 100 are more expensive at small $n$, but less expensive at large $n$.

## References

1. M.Abramowitz and I.A.Stegun, 'Handbook of Mathematical Functions', Dover (New York) 1965.
2. C.W.Clenshaw, 'N.P.L. Mathematical Tables', Vol. 5, HMSO (1962).
