

HSL ARCHIVE

1 SUMMARY

To compute the modified Bessel functions *Ibessnu* and Kv for a range of orders v=0,1,2,...,n, at the same argument x.

ATTRIBUTES — Version: 1.0.0. Types: FF07A; FF07AD. Calls: FD05A. Original date: April 1983. Origin: A. R. Curtis, Harwell.

2 HOW TO USE THE PACKAGE

2.1 Argument list and calling sequence

The single precision version

CALL FF07A(X,N,A,B,NAB)

The double precision version

CALL FF07AD(X,N,A,B,NAB)

- X is a REAL (DOUBLE PRECISION in the D version) variable which must be set by the user to the value of $\pm x$ (the subroutine uses the absolute value of X). A zero value is not allowed when $K_v(x)$ is requested, i.e. when argument N<0. This argument is not altered by the subroutine.
- N is an INTEGER variable which must be set by the user to $\pm n$, where *n* is the highest order required. If N<0, both $I_{\nu}(x)$ and $K_{\nu}(x)$ are computed; if N>0, only $I_{\nu}(x)$ is computed. A zero value is not allowed. This argument is not altered by the subroutine.
- A is a REAL (DOUBLE PRECISION in the D version) array of dimension NAB in which the subroutine will return the values of $I_v(x)$, v=0,1,2,...,n, i.e. $I_0(x)$ in A(1), $I_1(x)$ in A(2) up to $I_n(x)$ in A(n+1).
- B is a REAL (DOUBLE PRECISION in the D version) array of dimension NAB in which the subroutine will return the values of $K_v(x)$, v=0,1,2,...,n, i.e. $K_0(x)$ in B(1), $K_1(x)$ in B(2) up to $K_n(x)$ in B(n+1). If N>0 the array B is not altered.
- NAB is an INTEGER variable which must be set by the user to the dimension of the arrays A and B; it must be at least n+1. This argument is not altered by the subroutine.

3 GENERAL INFORMATION

Use of common: none.

Workspace: None.

Other routines called directly: calls DEXP, DSQRT and DLOG.

Input/output: Error warning and diagnostic messages on unit 6.

Restrictions:

n > 0;

x > 0 if $K_v(x)$ required,

else $x \ge 0$.

<u>FF07</u>

4 METHOD

4.1 Method

The method is based on the recurrence relation (9.6.26 in Abramowitz and Stegun^[1])

$$u_{\nu+1} - u_{\nu-1} = \frac{2\nu}{x} u_{\nu} \tag{1}$$

whose general solution is

$$u_{v} = a I_{v}(x) + b(-1)^{v} K_{v}(x)$$
⁽²⁾

where *a* and *b* are constants.

To compute $I_{\nu}(x)$, the recurrence relation is solved for $u_{\nu-1}$ and used downwards from $\nu=m>n$, chosen so that the second term in (2) is negligible compared with the first for $\nu \le n$. At the same time, the sum

$$s = u_0 + 2\sum_{\nu=1}^m u_{\nu} \approx a \, e^x \tag{3}$$

is computed, and the normalizing factor $a^{-1} = s^{-1}e^x$ is finally applied to the stored values in the array A.

To compute $K_{\nu}(x)$ (if required), the recurrence relation is solved for $u_{\nu+1}$ and used upwards from accurate values of $K_0(x)$ and $K_1(x)$. The first is computed from new highly accurate Chebyshev series, and $K_1(x)$ is then computed as (9.6.15^[1])

$$K_1(x) = \frac{1/x - I_1(x)K_0(x)}{I_0(x)},$$
(4)

which does not suffer serious loss of accuracy through cancellation.

4.2 Accuracy and timing

The computation of $I_0(x)$ is good almost to full computer accuracy. The values of $I_v(x)$ in Table 9.11 in Abramowitz and Stegun^[1] are reproduced exactly (apart from those which underflow). It may be necessary to take precautions against overflow. For small x and large n, underflows will occur and zero values will be returned; a call could be made to a library subroutine to mask off underflow interrupts.

The accuracy of $K_{\nu}(x)$ is also about 15 significant figures using the new Chebyshev series for $K_0(x)$. All $K_{\nu}(x)$ values in Table 9.11^[1] for $x \le 10$ are reproduced (except for a few discrepancies by one in the last digit of the table values, and for values which would overflow); however, for x = 50 and x = 100, many errors in the last digit of the table values where found, and the accuracy of FF07AD was confirmed independently by using the asymptotic series 9.7.2^[1]. For small *x* and large *n*, overflows would occur on the IBM computer; these are avoided, and the largest value represented on the machine is returned instead.

A 4-byte arithmetic version on some computers is not recommended because of possible accumulation of rounding errors; since the result arrays A and B are used as work-space, this consideration prevents internal 8-byte arithmetic computation being used to give 4-byte precision results. If FF07AD is used on other computers, the constants used in accuracy and overflow tests should be reconsidered.

Execution time depends more or less linearly on *n*, and also somewhat on *x*. For moderate *x* values, it varies on an IBM/3081 from about 75 μ sec at *n* = 10 to about 650 μ sec at *n* = 100 for $I_v(x)$ only; if $K_v(x)$ is also computed, typical times are 125 μ sec and 800 μ sec. Values of *x* in the range 10 to 100 are more expensive at small *n*, but less expensive at large *n*.

References

- 1. M.Abramowitz and I.A.Stegun, 'Handbook of Mathematical Functions', Dover (New York) 1965.
- 2. C.W.Clenshaw, 'N.P.L. Mathematical Tables', Vol. 5, HMSO (1962).